

Sieve weights and their smoothings

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Selberg's sieve

$$\sum_{d|(a,m)} \mu(d) \leq \left(\sum_{d|(a,m)} \lambda_d \right)^2,$$

for any $\lambda_d \in \mathbb{R}$ with $\lambda_1 = 1$.

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$$\sum_{a \in \mathcal{A}} \left(\sum_{d|(a,m)} \lambda_d \right)^2 = \sum_{\substack{d_1, d_2 | m \\ D = [d_1, d_2]}} \lambda_{d_1} \lambda_{d_2} \cdot \#\{a \in \mathcal{A} : D|a\}.$$

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Optimizing (making assumptions on \mathcal{A}) yields

$$\lambda_d \approx c \cdot \mu(d) \cdot \left(\frac{\log(R/d)}{\log R} \right)^\kappa \cdot \mathbf{1}_{d \leq R}, \quad \kappa = \textit{sieve dimension}$$

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$\text{supp}(f) \subset (-\infty, 1]$, $f(0) = 1$.

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$M_f(n; R)$ should behave like a sieve weight for f sufficiently smooth, i.e.

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[Maynard](#) and [Tao](#) used k -dimensional generalization of $M_f(n; R)$ to detect small gaps between primes.

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If $n = 2m$, $2 \nmid m$, then

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Ford: $\#\{n \leq x : \exists! d \in (R/2, R], d|n\} \asymp x(\log R)^{-\delta}(\log \log R)^{-3/2}$,

where $\delta \approx 0.086 < 1$, i.e. $\sum_{d|m, d \leq R} \mu(d) \neq 0$ too often.

How much to smooth?

For $f \in C^1(\mathbb{R})$, $n = p^v m$ with $p \nmid m$,

$$M_f(n; R) = \sum_{d|m} \mu(d) f\left(\frac{\log d}{\log R}\right) + \sum_{d|m} \mu(pd) f\left(\frac{\log(pd)}{\log R}\right)$$

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For $f \in C^A(\mathbb{R})$, $n = p_1^{v_1} \cdots p_A^{v_A} m$ with $p_1, \dots, p_A \nmid m$,

$$M_f(n; R) = (-1)^A \int_0^{\frac{\log p_1}{\log R}} \cdots \int_0^{\frac{\log p_A}{\log R}} \sum_{d|m} \mu(d) f^{(A)}\left(\sum_{a=1}^A u_a + \frac{\log d}{\log R}\right) d\mathbf{u}$$

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Guess :
$$\sum_{n \leq x} M_f(n; R)^{2k} \lesssim \max \left\{ \frac{x}{\log R}, \frac{\sum_{n \leq x} (\sum_{d|n, d \leq R} \mu(d))^{2k}}{(\log R)^{2kA}} \right\}$$

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If $\sum_{n \leq x} (\sum_{d|n, d \leq R} \mu(d))^{2k} \sim c_k x (\log R)^{E_k}$, then we would need $A > E_k/2k$ for $M_f(n; R)^{2k}$ to act as a sieve weight.

Theorem (Granville, K., Maynard (201?))

Let $k, A \in \mathbb{N}$. Assume that:

- $\sum_{n \leq x} (\sum_{d|n, d \leq R} \mu(d))^{2k} \sim c_k x (\log R)^{E_k}$ when $x/R^{2k} \rightarrow \infty$;
- $f \in C^A(\mathbb{R})$, $\text{supp}(f) \subset (-\infty, 1]$, $f, f', \dots, f^{(A)}$ unif. bounded.

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(a) If $A > E_k/2k + 1$, then

- $$\sum_{\substack{n \leq x \\ \exists p|n, p \leq R^\eta}} M_f(n; R)^{2k} \ll_{f,k,A} \frac{\eta^{3/2} x}{\log R} \quad (x \geq R \geq 2)$$

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- $$\sum_{n \leq x} M_f(n; R)^{2k} = \frac{c_{k,f} x}{\log R} + O_{f,k,A} \left(\frac{x}{(\log R)^{3/2}} \right) \quad (x \geq R^{2k+\epsilon})$$

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(b) If $A \leq E_k/2k + 1$ and $x \geq R \geq 2$, then

$$\sum_{n \leq x} M_f(n; R)^{2k} \ll x (\log R)^{E_k - 2k(A-1)}.$$

Dominant contribution when $\#\{p|n : p \leq R\} \sim (E_k + 2k) \log \log R$.

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Finite field analogy:

$$\frac{1}{q^n} \sum_{\substack{F \in \mathbb{F}_q[t] \\ \deg(F) = n}} \left(\sum_{G|F, \deg(G) = m} \mu(G) \right)^{2k}.$$

The analogy for permutations

$$\text{Perm}(n; m, k) := \frac{1}{n!} \sum_{\sigma \in \mathcal{S}_n} \left(\sum_{\substack{T \subset [n], |T|=m \\ \sigma(T)=T}} \mu(\sigma|_T) \right)^{2k}$$

with $\mu(\tau) := (-1)^{\#\{\text{cycles of } \tau\}}$.

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We have $2^{2k-1} - 1$ free variables of size $\leq m$ and $2k$ constraints :

$$\text{Perm}(n; m, k) \asymp m^{2^{2k-1} - 2k - 1} \quad (k \geq 2).$$

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Guess: $E_k = 2^{2k-1} - 2k - 1$ in $\sum_{n \leq x} \left(\sum_{d|n, d \leq R} \mu(d) \right)^{2k} \sim c_k x (\log R)^{E_k}$.

The analogy for permutations, 3

$$\text{Perm}(n; m, k) = \#\{(r_l)_{\emptyset \neq I \subset [2k]} : r_l \in \{0, 1\} \text{ (}|I| \text{ odd)}, \sum_{l: i \in I} r_l = m, \forall i\}.$$

We have $2^{2k-1} - 1$ free variables of size $\leq m$ and $2k$ constraints :

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Motohashi proved $E_2 = 2$, but $2^{4-1} - 4 - 1 = 3$ (the analogy between \mathbb{Z} and $\mathbb{F}_q[t]$ analogy breaks down).

The source of the extra cancellation

Write $d \approx R$ if $R/2 < d \leq R$, so that

$$\frac{1}{x} \sum_{n \leq x} \left(\sum_{d|n, d \approx R} \mu(d) \right)^{2k} \sim \sum_{d_1, \dots, d_{2k} \approx R} \frac{\mu(d_1) \dots \mu(d_{2k})}{[d_1, \dots, d_{2k}]}.$$

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The total main term cancels since $\sum_n \mu(n)/n = 0$.

$$\implies E_k < 2^{2k-1} - 2k - 1.$$

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$$\sum_{n \leq x} \left(\sum_{d|n, d \leq R} \mu(d) \right)^{2k} \sim \frac{x}{(2\pi i)^{2k}} \int \cdots \int_{\substack{\Re(s_j) = \frac{1}{\log R} \\ 1 \leq j \leq 2k}} F(\mathbf{s}) \frac{\prod_l^{\text{even}} \zeta(1 + s_l)}{\prod_l^{\text{odd}} \zeta(1 + s_l)} \prod_{j=1}^{2k} \frac{R^{s_j}}{s_j} d\mathbf{s},$$

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Theorem (Granville, K., Maynard (201?))

$$E_k = \binom{2k}{k} - 2k, \quad \text{i.e.} \quad \sum_{n \leq x} \left(\sum_{d|n, d \leq R} \mu(d) \right)^{2k} \sim c_k x (\log R)^{\binom{2k}{k} - 2k}.$$

Interpolating between integers and polynomials

If $\sigma \in \mathcal{S}_m$, then $\mu(\sigma) = (-1)^m \text{sgn}(\sigma)$ and $\mathbb{E}_{\rho \text{ cycle}}[\text{sgn}(\rho)] = 0$.

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(b) If $L(\beta, \chi) = 0$ and $(\log q)^{C_k} \leq \log R \leq \frac{1}{1-\beta}$, then

$$\frac{1}{x} \sum_{n \leq x} \left(\sum_{d|n, d \approx R} \chi(d) \right)^{2k} = (\log q)^{O_k(1)} \cdot (\log R)^{\binom{2k}{k} - 2k}.$$

Thank you!

$$\begin{aligned}
\text{Poly}_q(n, m; k) &:= \frac{1}{q^n} \sum_{\substack{N \in \mathbb{F}_q[t] \\ \deg N = n}} \left| \sum_{\substack{M|N \\ \deg M = m}} \mu(M) \right|^{2k} \\
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where

$$\mathcal{Z}_q(w) = \sum_{G \in \mathbb{F}_q[t]} \left(\frac{w}{q} \right)^{\deg(G)} = \prod_P (1 - (w/q)^{\deg(P)})^{-1},$$

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The torsion of \mathbb{R}/\mathbb{Z} , i.e. the discrete structure of $\mathbb{F}_q[t]$, yields a fundamentally different pole structure.